




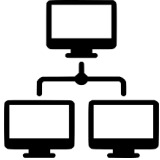
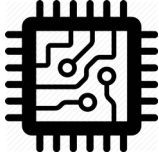
Data Processing at the Speed of 100 Gbps@Apache Crail (Incubating)

<http://crail.incubator.apache.org/>

Animesh Trivedi
IBM Research, Zurich

The Performance Landscape

Trend 1: The I/O performance is increasing dramatically

	2010	2018	2018-
 Storage	100 MB/s (HDDs)	1,000 MB/s (SSDs, NVMe)	10,000 MB/sec (3DXP)
 Network	1 Gbps	10/25/40 Gbps	100/200 Gbps
 CPU	1 x ~3 GHz	n x ~3 GHz	n x ~3 GHz

The Performance Landscape

Trend 2: The I/O diversity is increasing dramatically



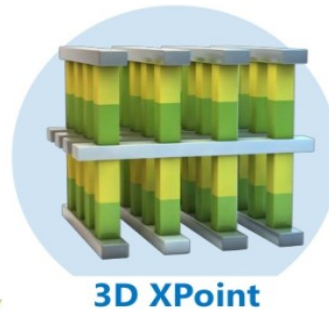
IBM Cloud



Google Cloud Platform

memory

data
latency
RDMA
system
network
operating
transfer
pinning
overhead



nvmTM
EXPRESS



**SPEED
LIMIT
100
Gbps**



The Key Challenge

Given the current hardware performance and diversity landscape, how can we orchestrate data movement at the speed of hardware

in other words:

"Can we feed modern data processing stacks at 100+ Gbps data speeds with 1-10 usec access latencies"

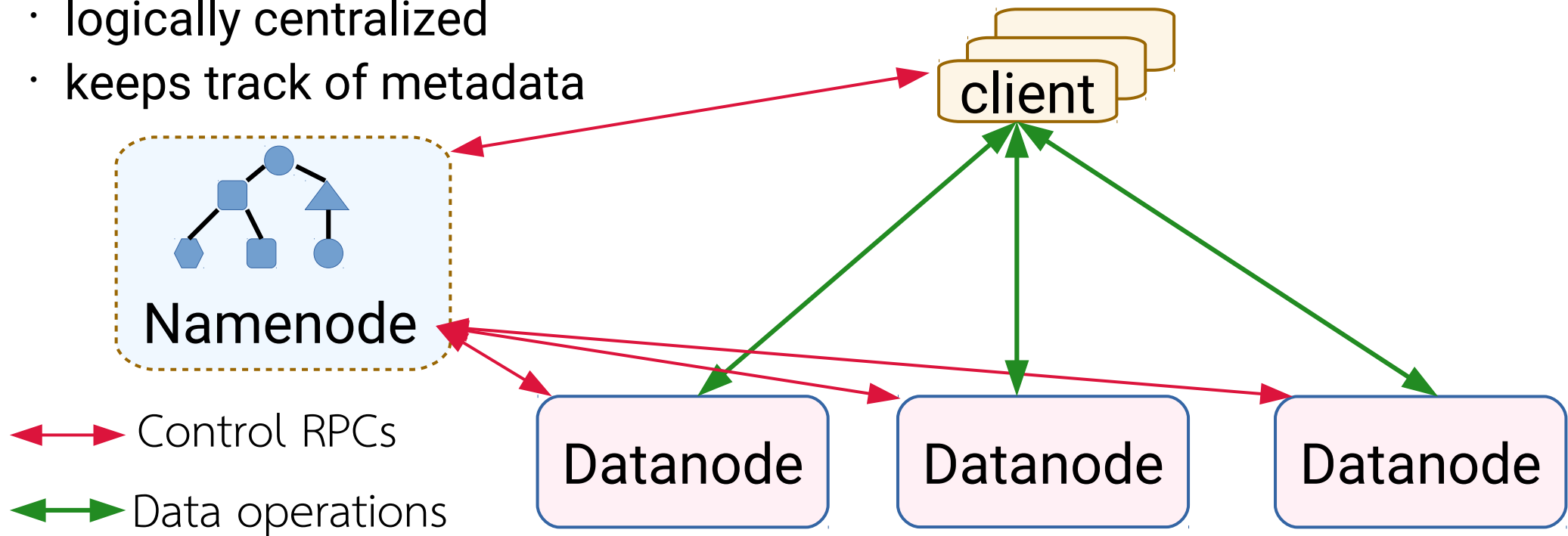
Apache Crail (Incubating)

- A distributed data “store” platform designed from scratch to “share” data at the speed of hardware
 - “store”: in DRAM/Flash/3DXP, with multiple front-end APIs to storage
 - “share”: intermediate data from intra-job (shuffle, broadcast) & inter-job
- Targets to speed-up workloads by accelerating data sharing
- An effort to unify isolated and incompatible efforts to integrate high-performance device in big data frameworks
 - Written in Java8 with multiple client-language support
- Started at IBM Research Lab, Zurich
- Apache Incubator project since November, 2017

System Overview – Crail Store

- logically centralized
- keeps track of metadata

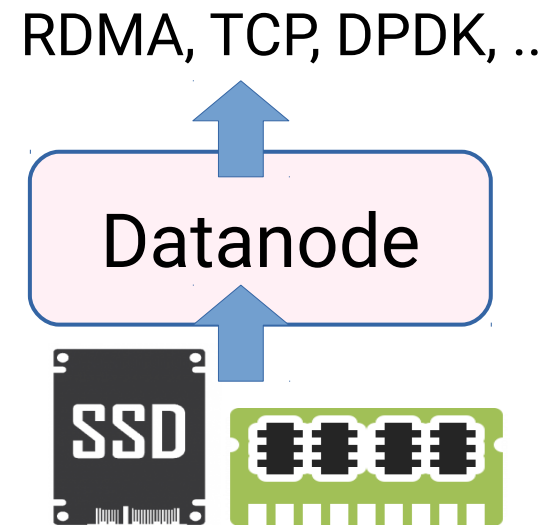
- access metadata from the namenode
- access data from datanodes



- donate storage and network resources to Crail
- mostly dumb/stateless servers

Apache Crail: Datanode

- DataNode is responsible for exposing (storage + networking) resources to the system
 - DRAM + RDMA [1], TCP/Sockets
 - Flash + NVMeF [2]
 - Local DRAM + memcpy
 - Local flash + SPDK
- Accepts client connections and serves data
- Periodic heartbeat pings with the namenode



Apache Crail: Namenode

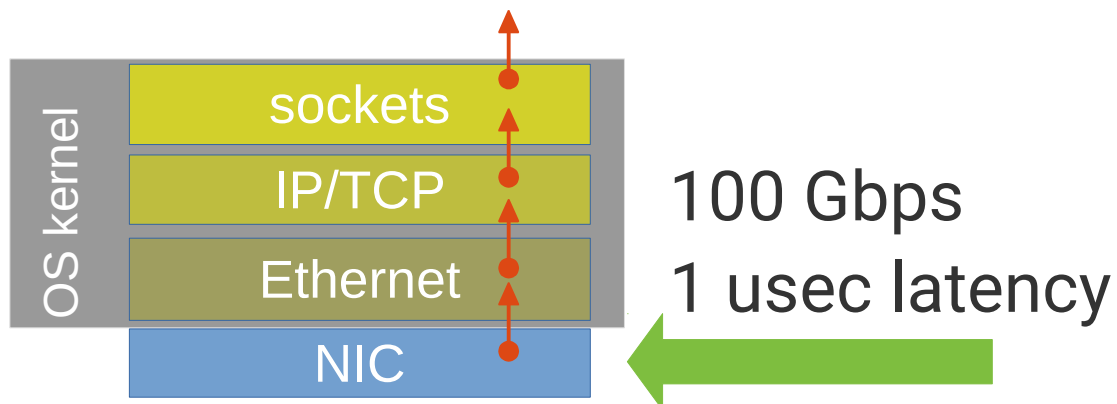
- Responsible for keeping track of storage resources in the cluster
 - Done by managing them in blocks (e.g., 1MB block size)
- Maintains a hierarchical node tree
 - directory, data files, stream files, multfiles, tables, KV
- Clients connect to the namenode to create new nodes, lookup, read, and write to nodes
 - Allocation policy (different media and node types)
 - A node can contains some blocks from DRAM and flash

Apache Crail: Clients

- Clients read and write data
 - Standalone or a part of a compute framework
 - Single writer, without holes files
 - Distributed clients often implicitly index and synchronize on the file/directory name
- Multiple-client side storage abstractions/interfaces
 - Simple hierarchical file system
 - Key-Value (KV) store
 - HDFS
 - Streaming (WiP)

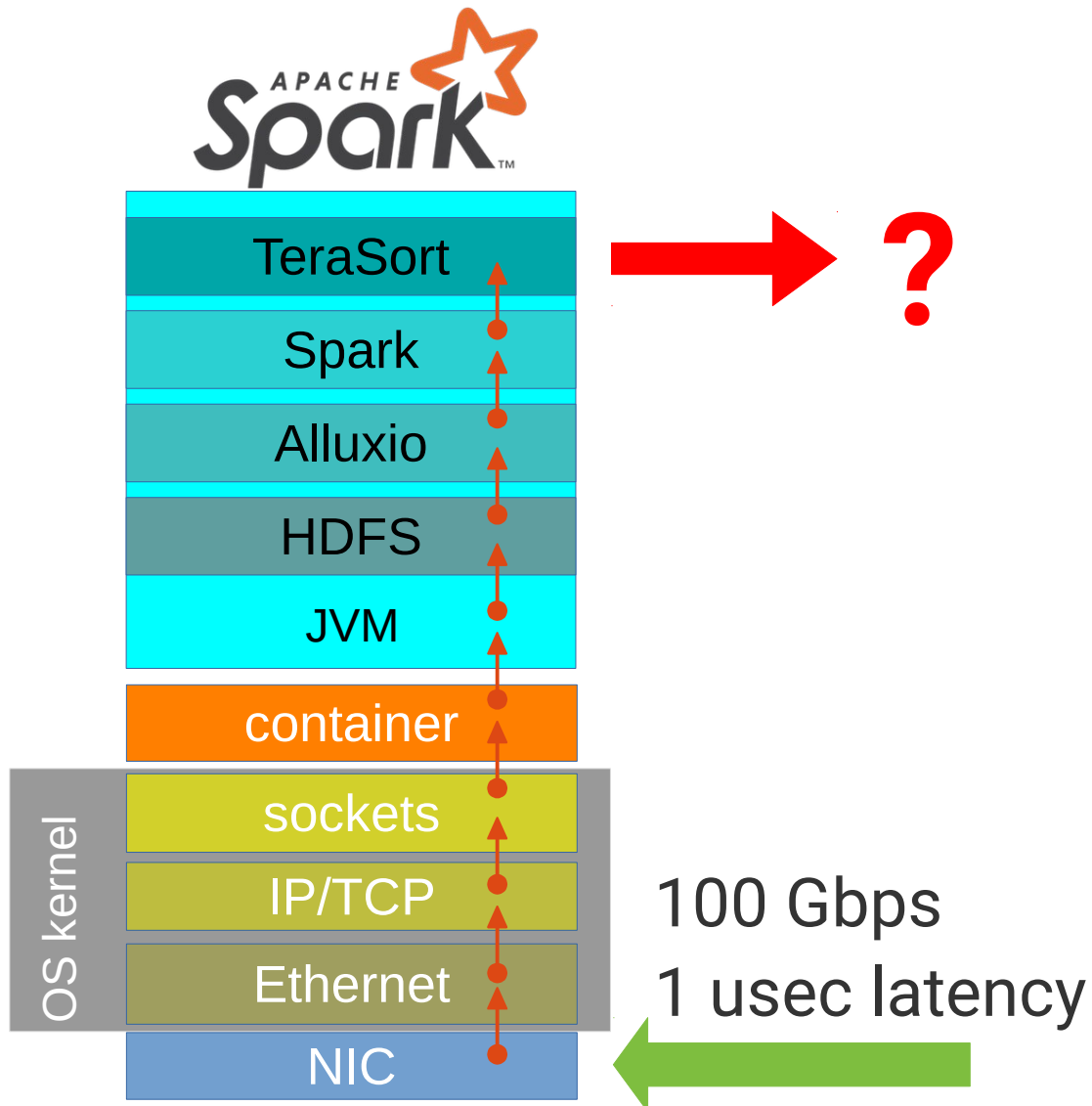
So where does
the performance come from?

Performance Principles – I



DataWorks Summit 2018, San Jose

Performance Principles – I

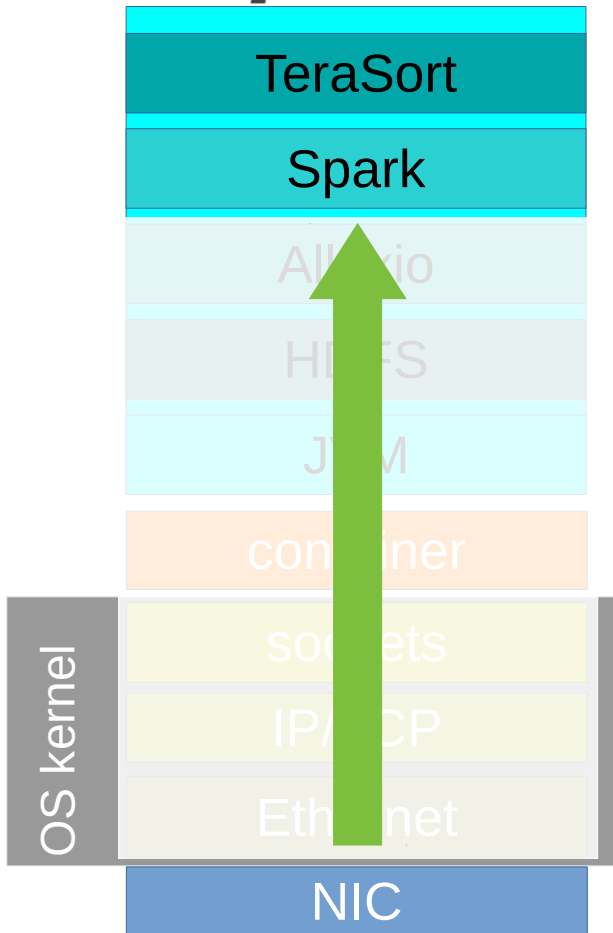


Performance Principles – I



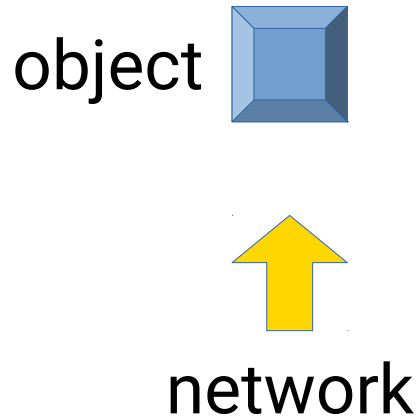
1. Use high-performance user-level I/O

RDMA[1], NVMeF[2], SPDK/DPDK in Java
- one-sided RDMA read/write operations

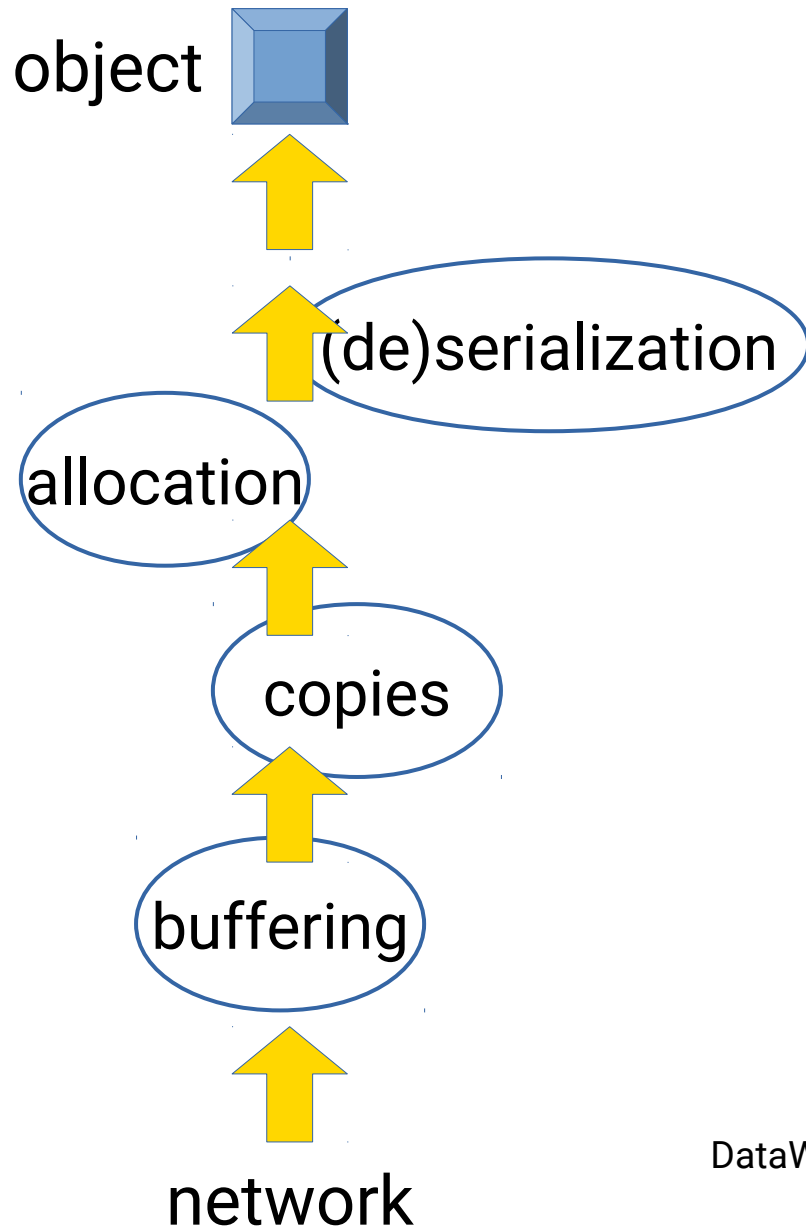


100 Gbps
1 usec latency

Performance Principles – II

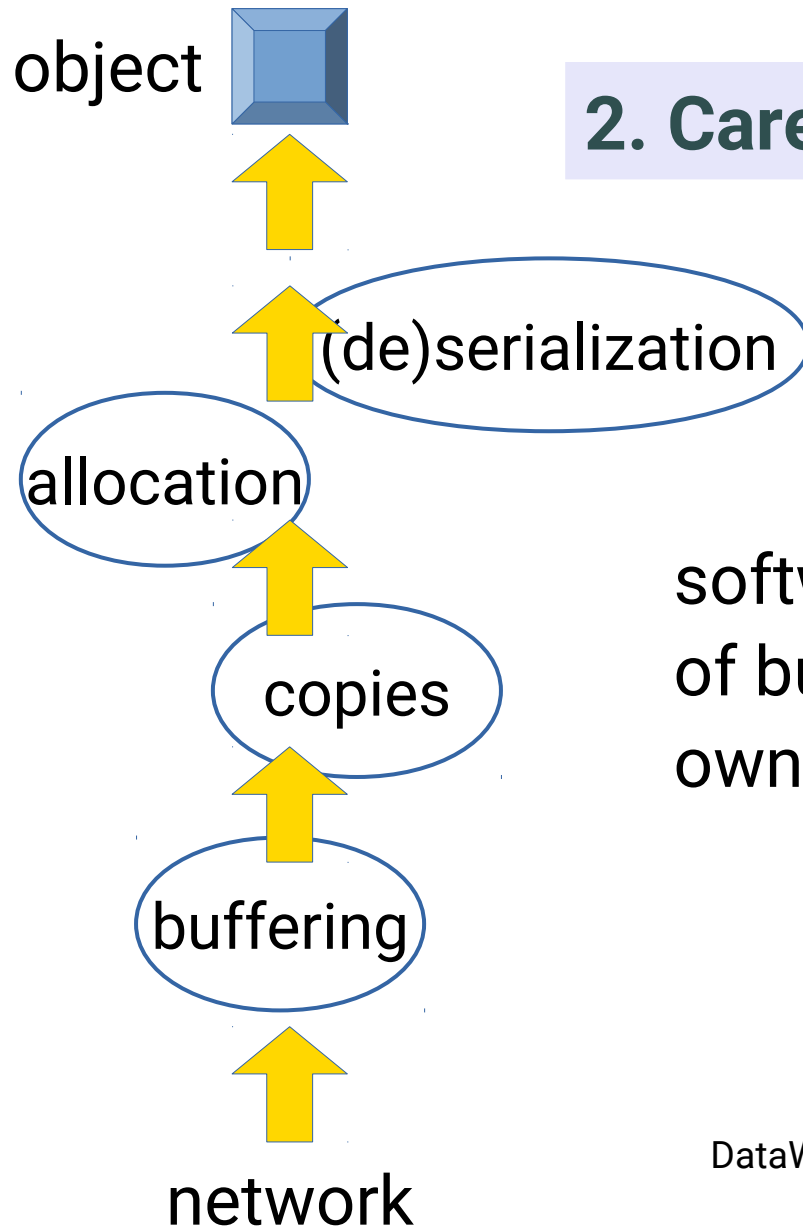


Performance Principles – II



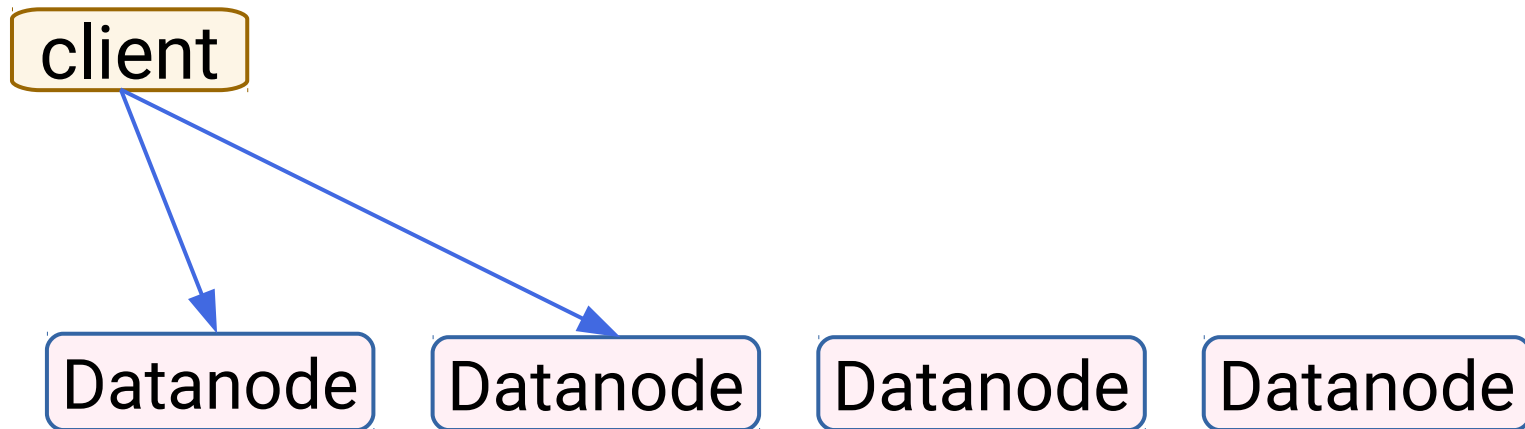
Performance Principles – II

2. Careful software design for the μ sec-Era

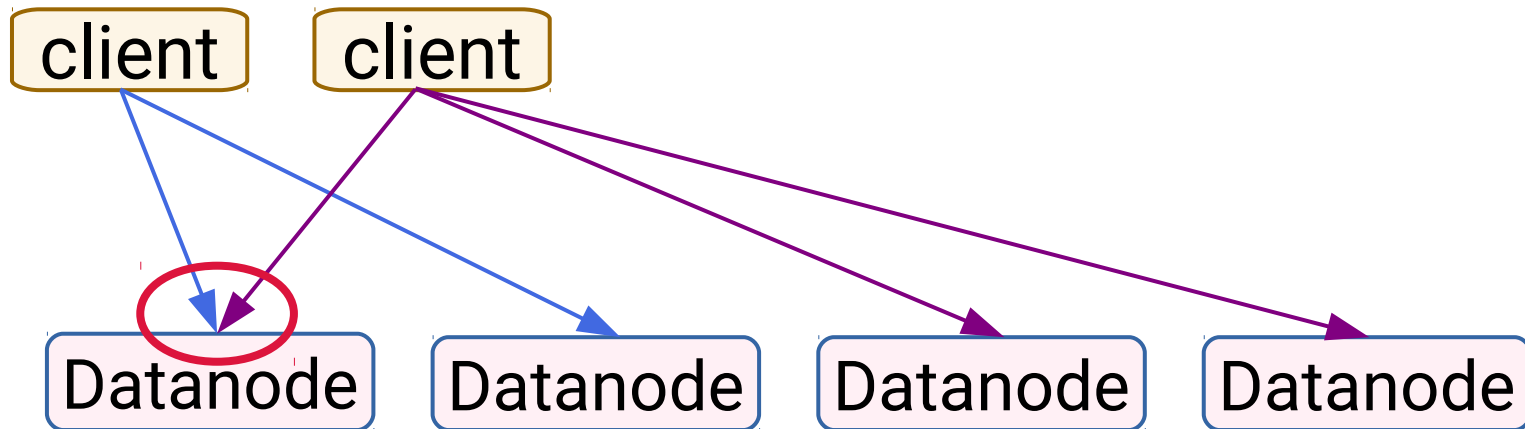


software time budgeting, pre-allocation of buffers, buffer caching, and reusing, ownership, etc.

Performance Principles – III



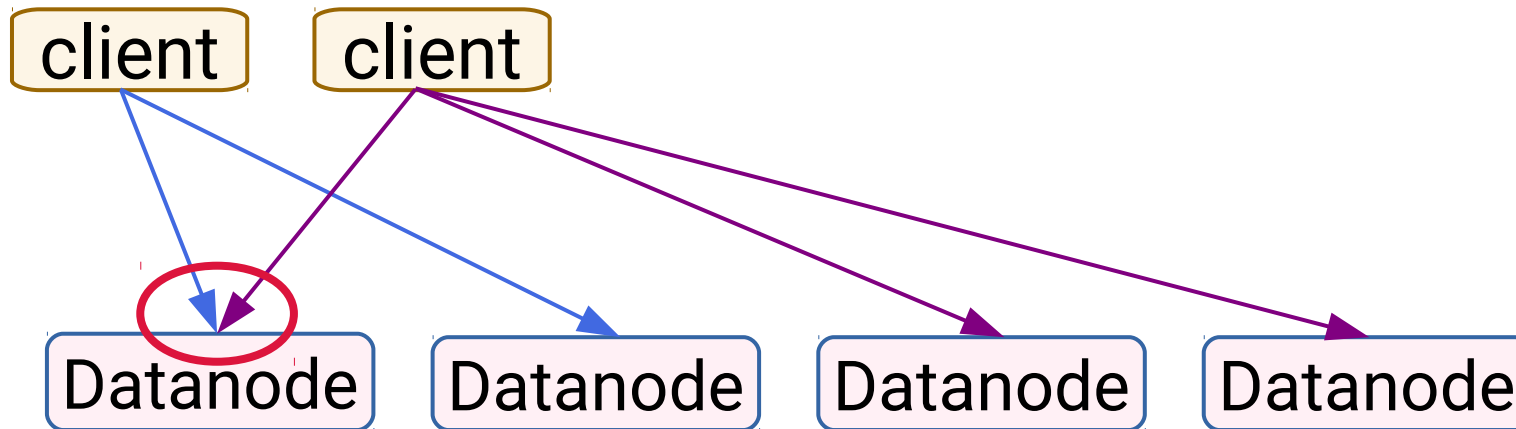
Performance Principles – III



Performance Principles – III

3. Careful data orchestration in the cluster

randomization, async requests, compute-I/O overlap, smart buffering, data allocation policies, etc.



Performance Principle - Recap

1. Use high-performance user-level I/O

2. Careful software design for the μ sec-Era

3. Careful data orchestration in the cluster

But what If I don't have ... RDMA

1. Use high-performance user-level I/O



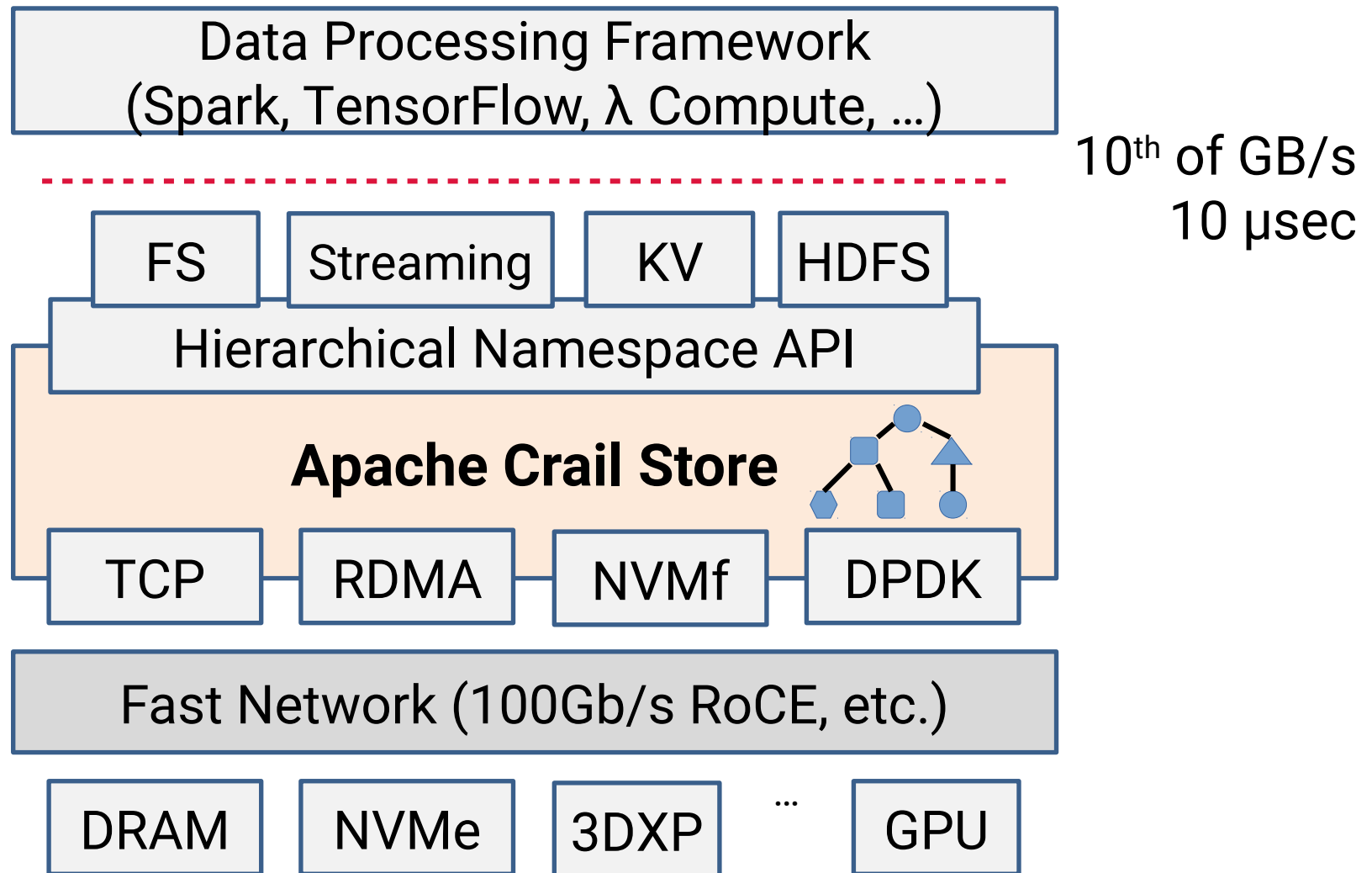
2. Careful software design for the μ sec-Era



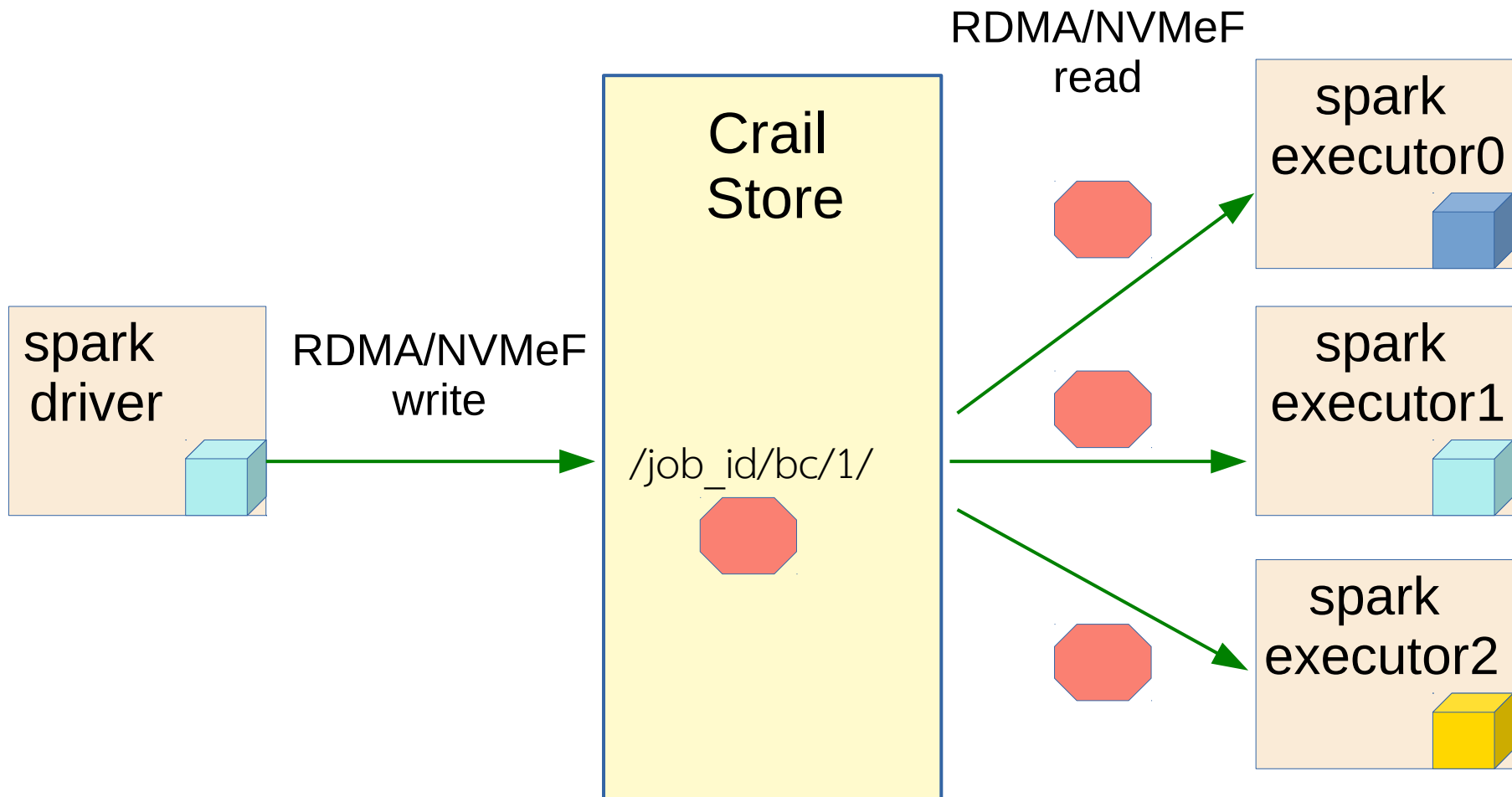
3. Careful data orchestration in the cluster



The Full Apache Crail (Incubating) Stack

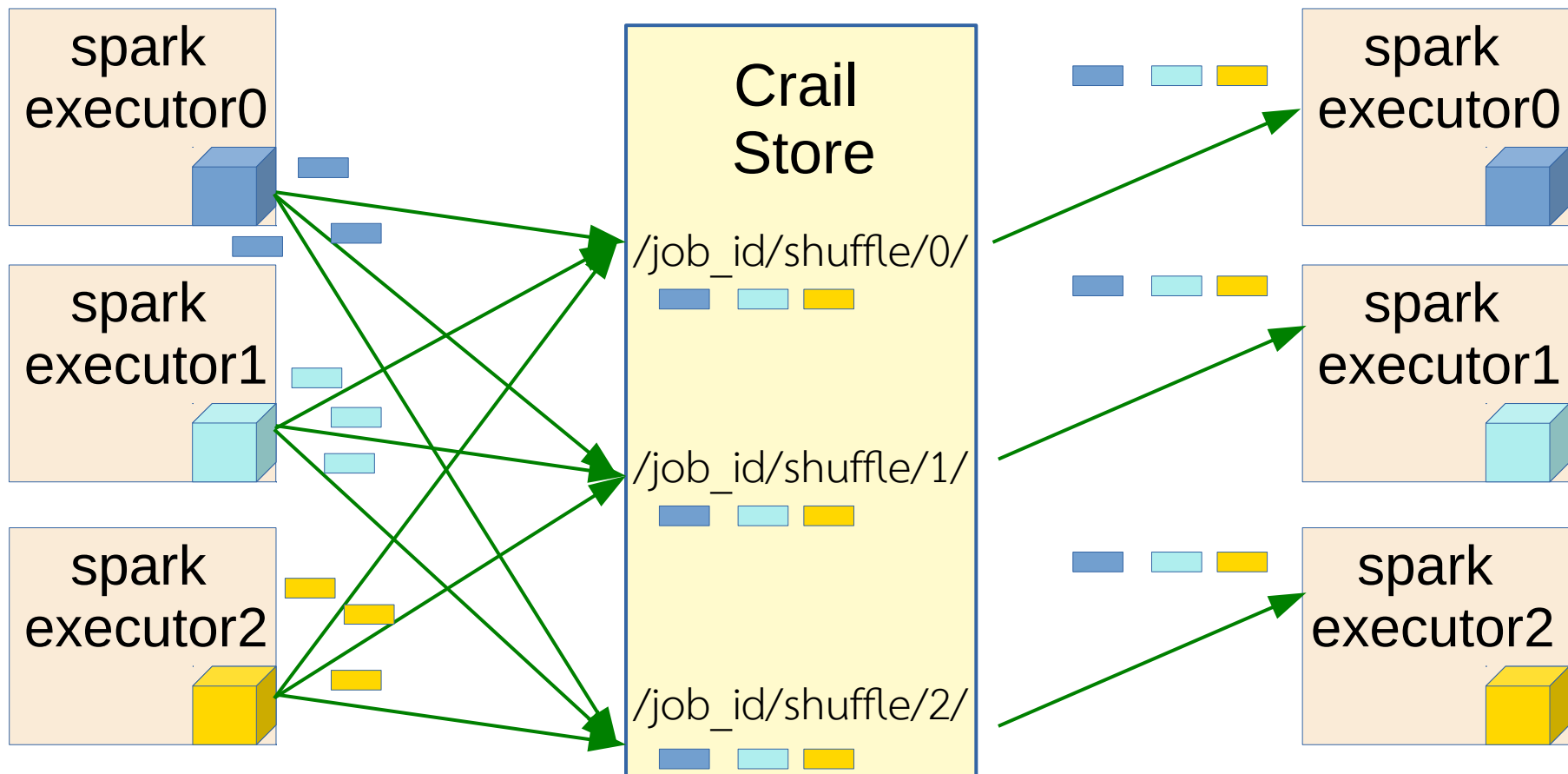


Use in Apache Spark – Broadcast



Use in Apache Spark - Shuffle

Spark Shuffle plugin [4] that write and reads data in Crail

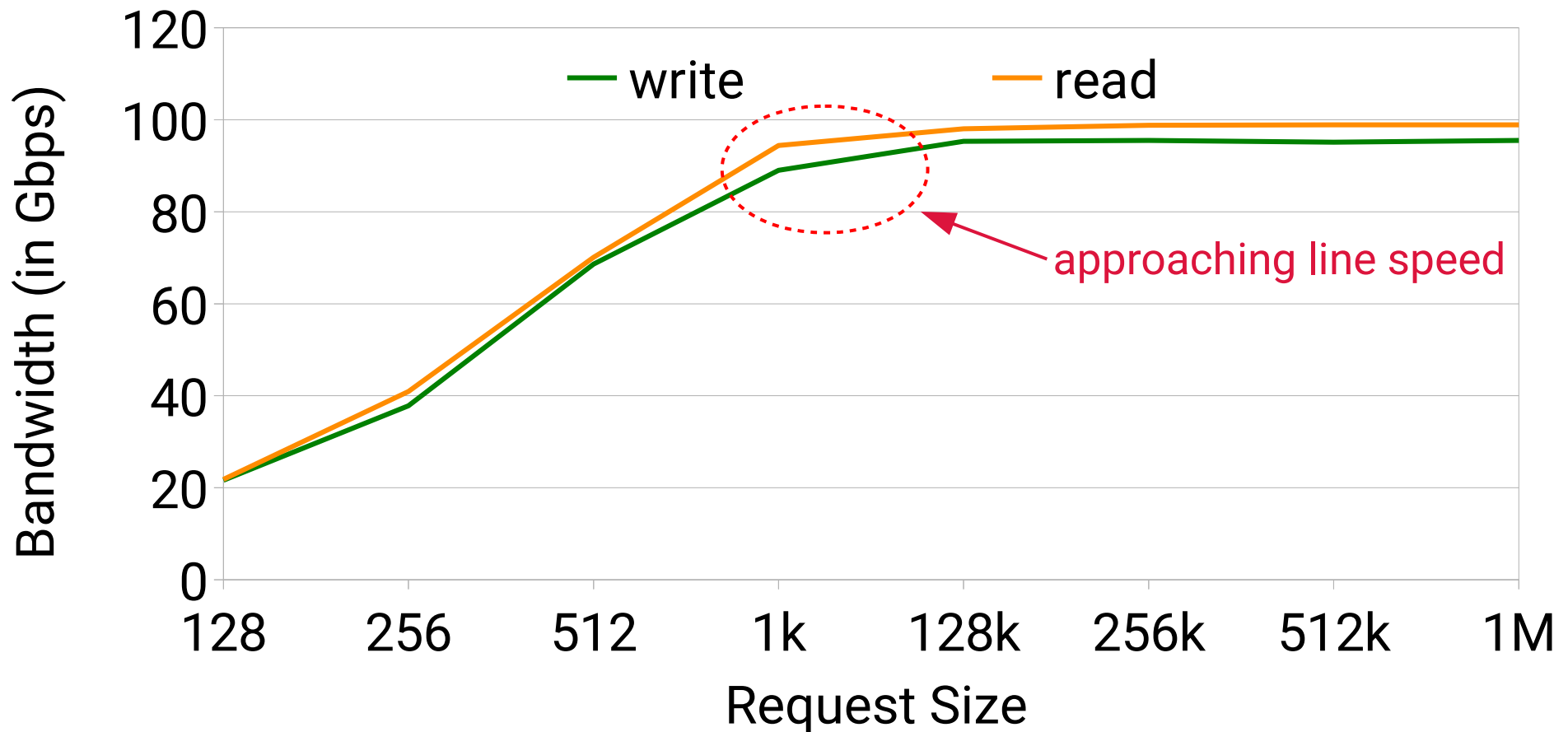


Performance Numbers

- Baseline performance for Apache Crail Store [6]:
 - Latency, bandwidth, and IOPS numbers in *a distributed setting in the JVM*
- Crail + Spark integration, available at [4]
 - Micro-operations: Broadcast and GroupBy
 - Workloads: TeraSort and SQL JOIN
 - Mixed DRAM and Flash setting
- A mix of x86 and POWER8 machines, 100 GbE network, 256GB DRAM DDR3, and NVMe flash

Crail Store – DRAM

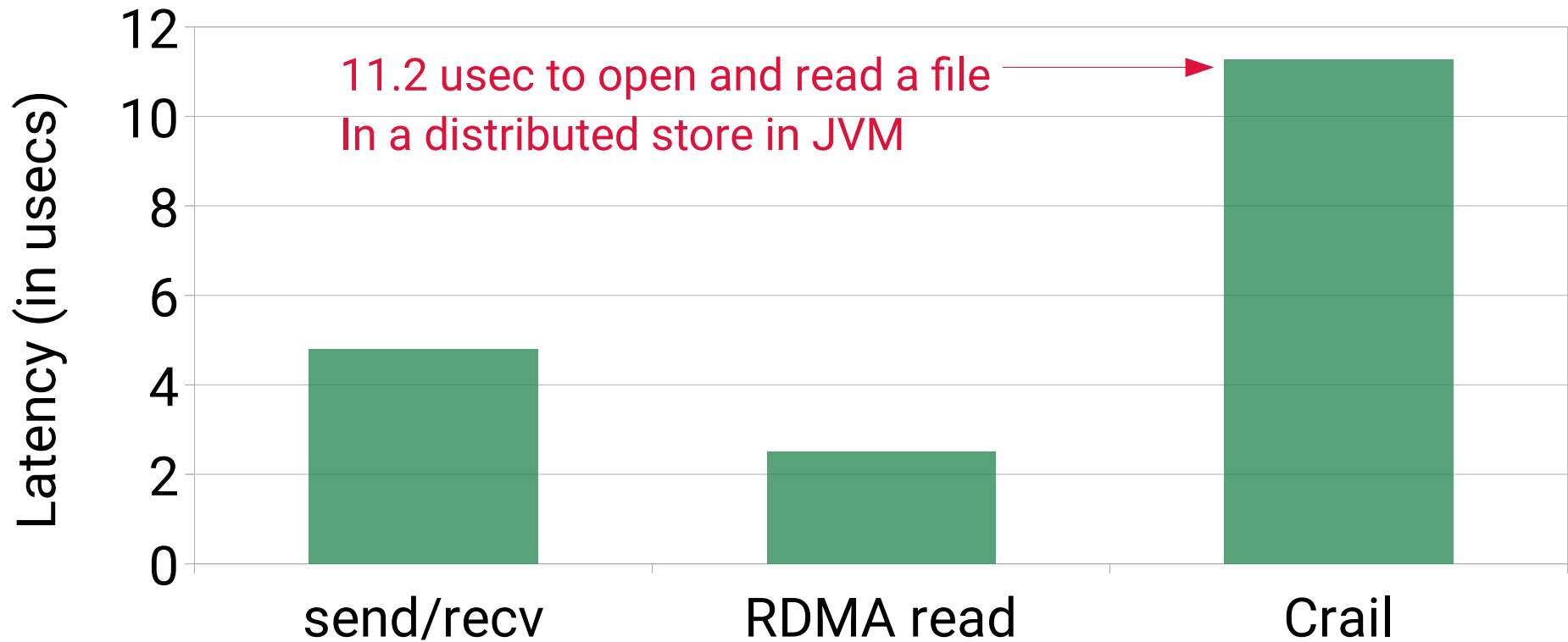
“File Read Bandwidth”



Crail delivers full hardware bandwidth from DRAM

Crail Store – DRAM

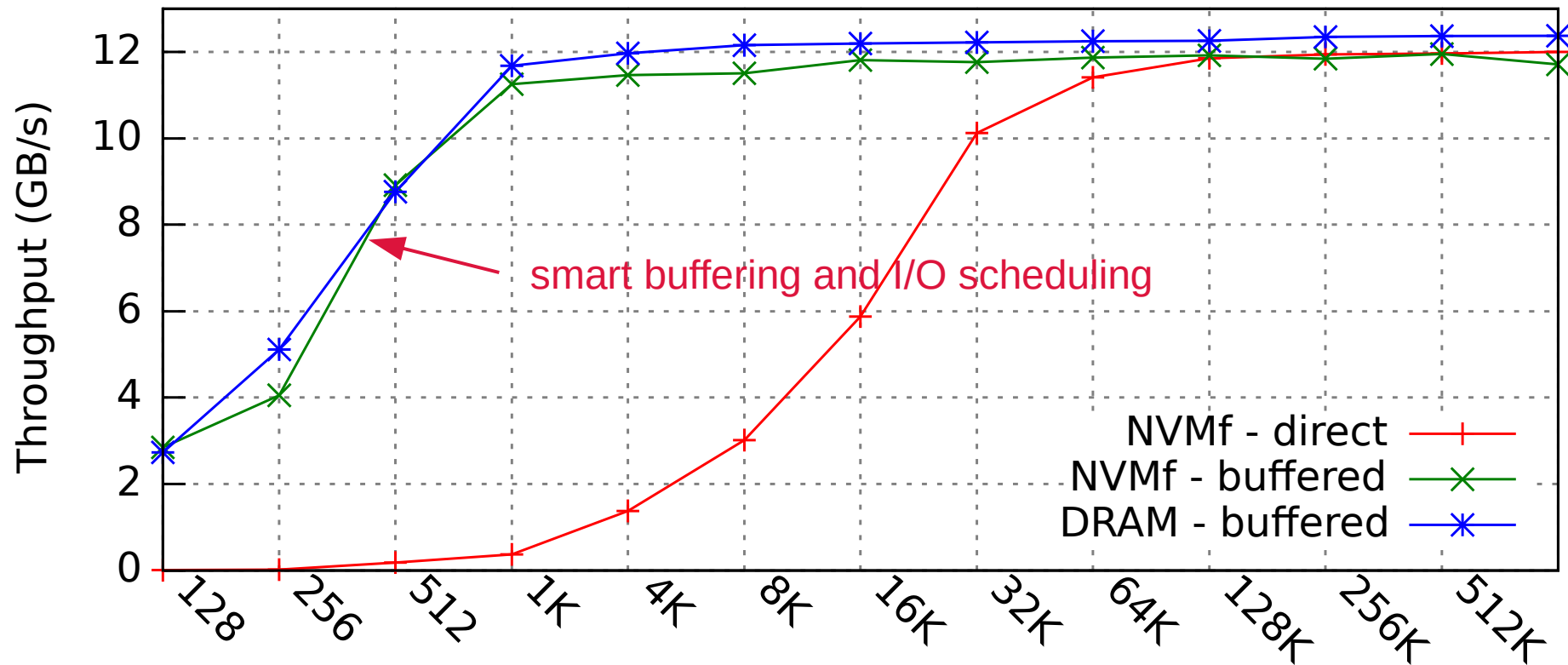
“File Read Latency”



Crail delivers ultra-low file/data access latencies in JVM

Crail Store – NVMeF

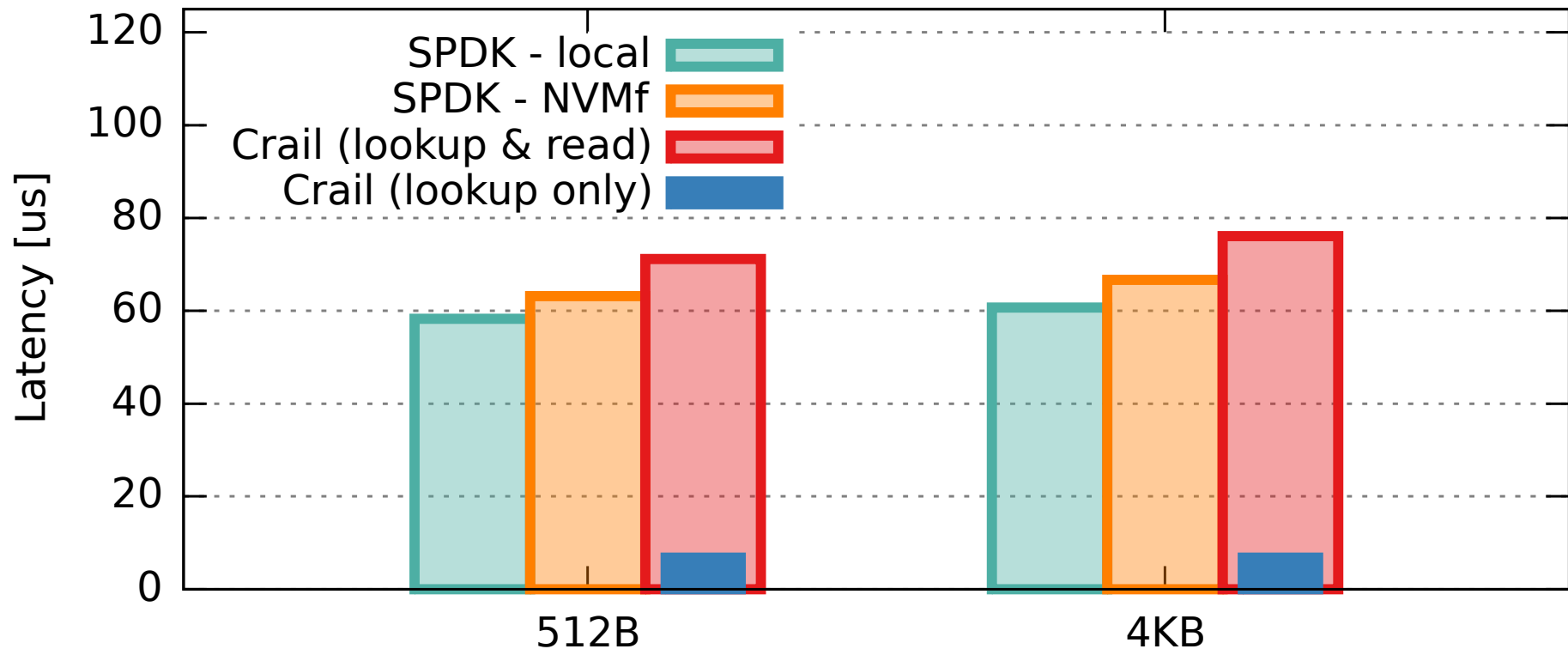
“File Read Bandwidth”



Crail delivers full hardware bandwidth from remote NVMe flash

Crail Store – NVMeF

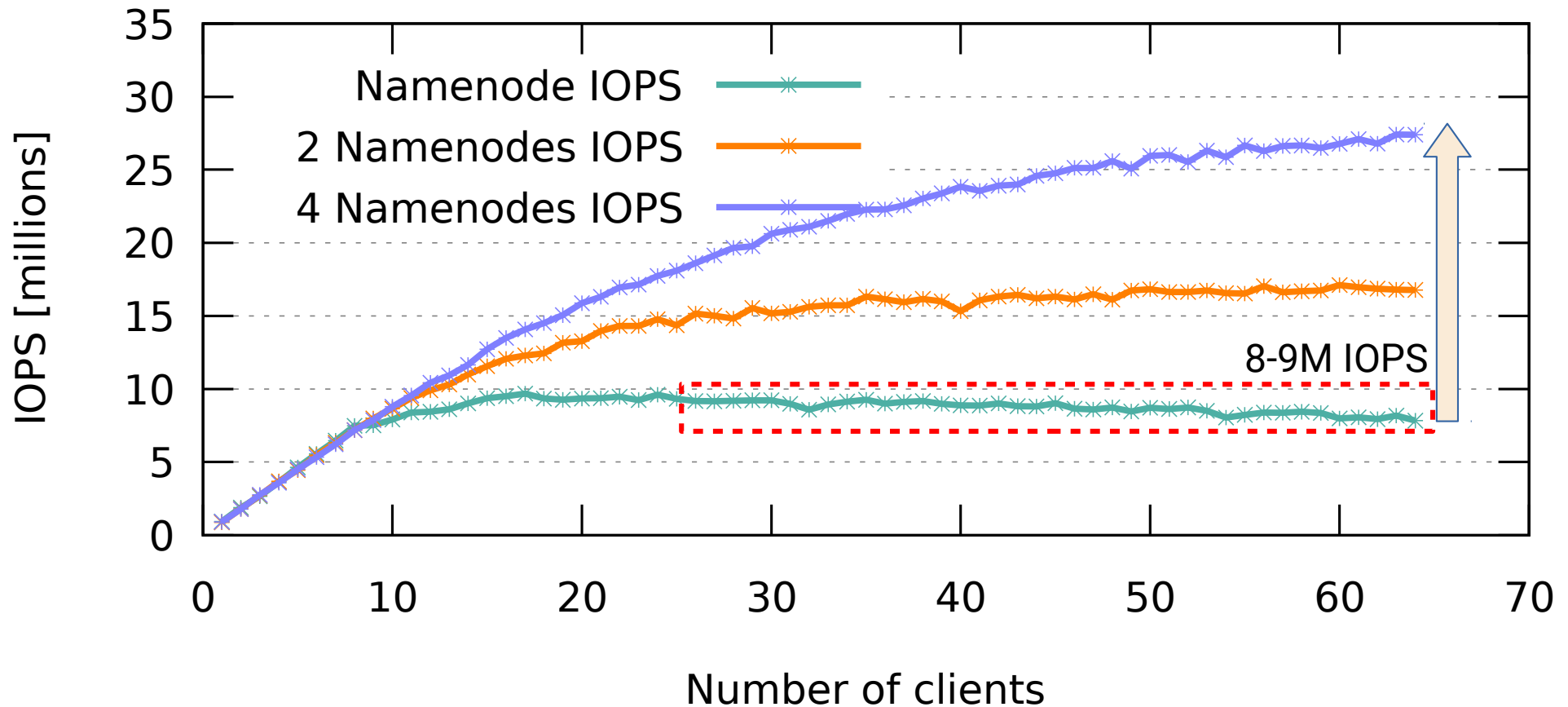
“File Read Latency”



Crail delivers native NVMe access performance in JVM

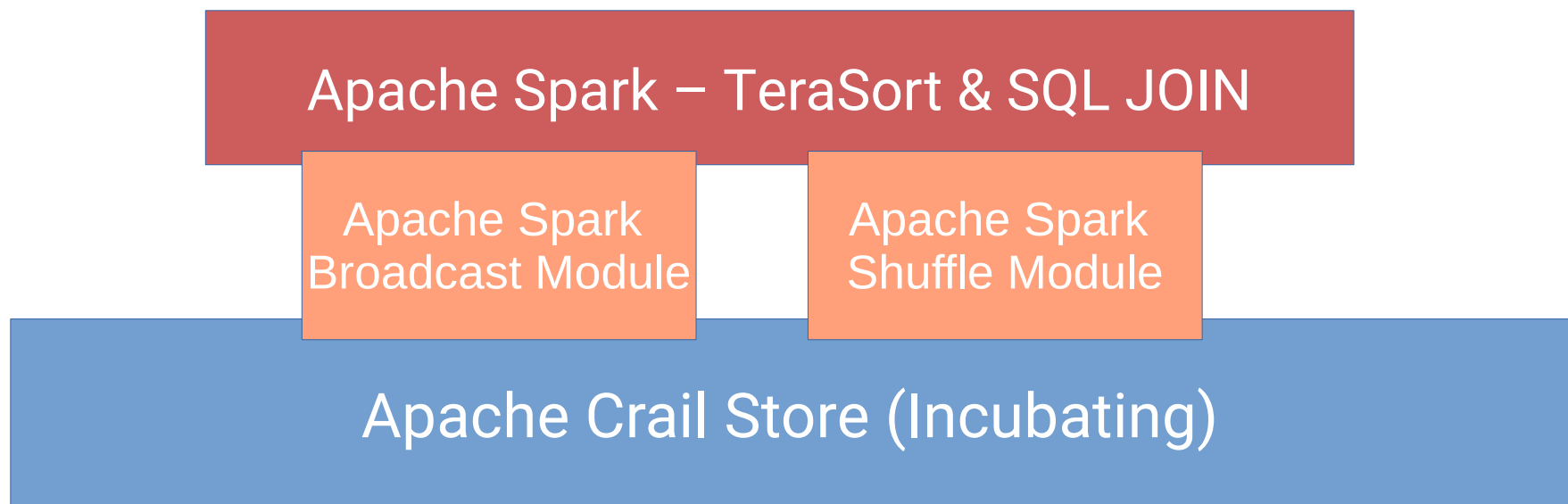
Crail Store – Metadata IOPS

“getFile operation”

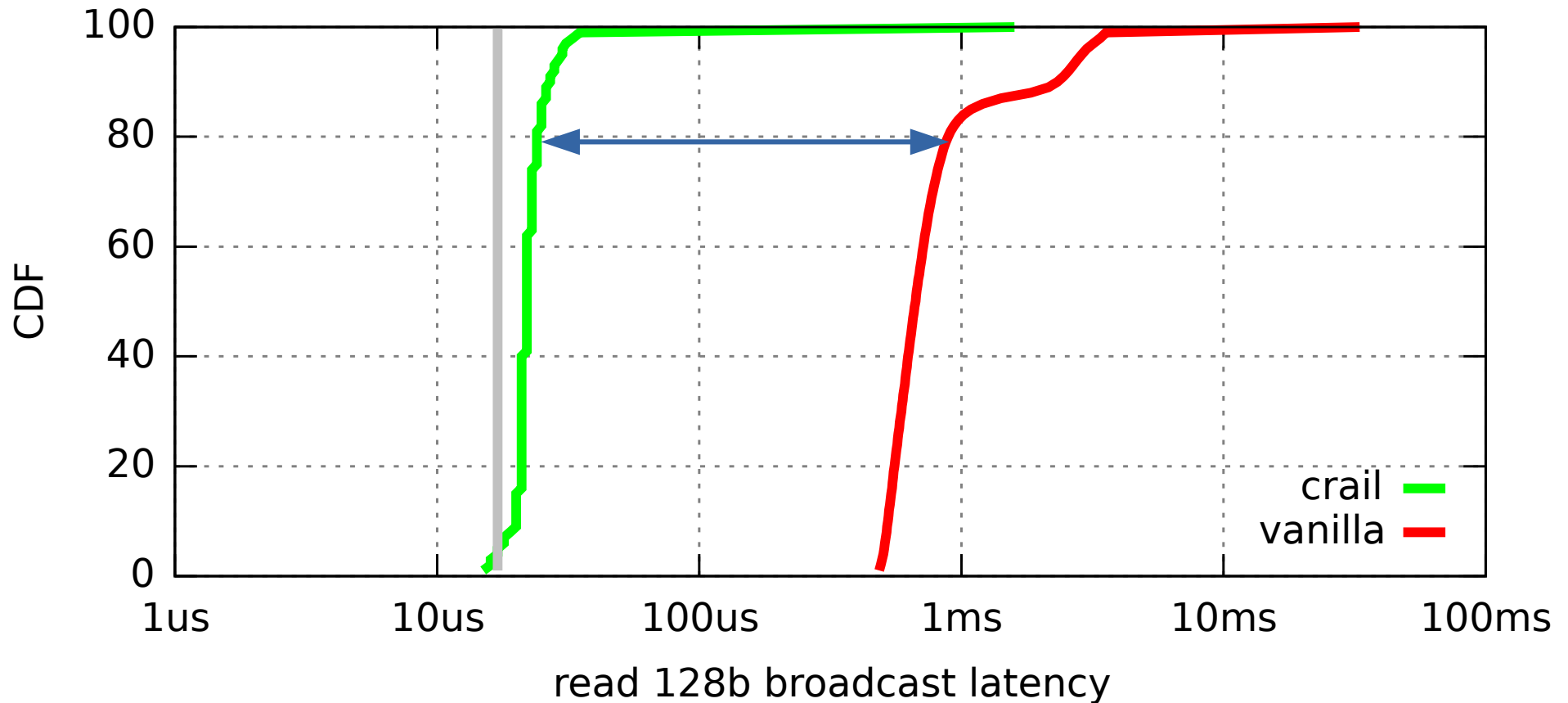


Crail offers (almost) linear scalability with Namenode ops

How does data processing-level performance look like?

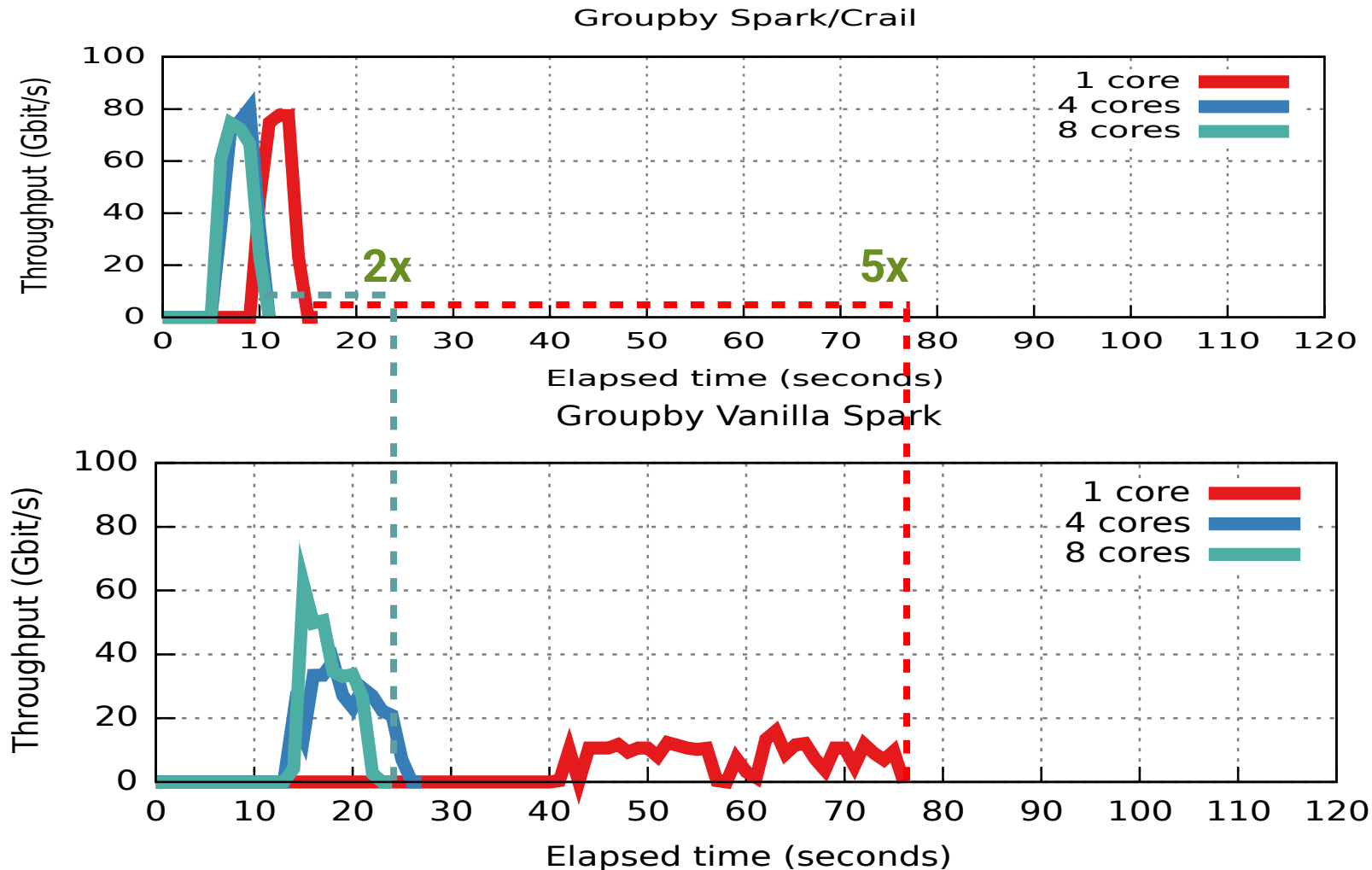


Crail Workload - Spark - Broadcast



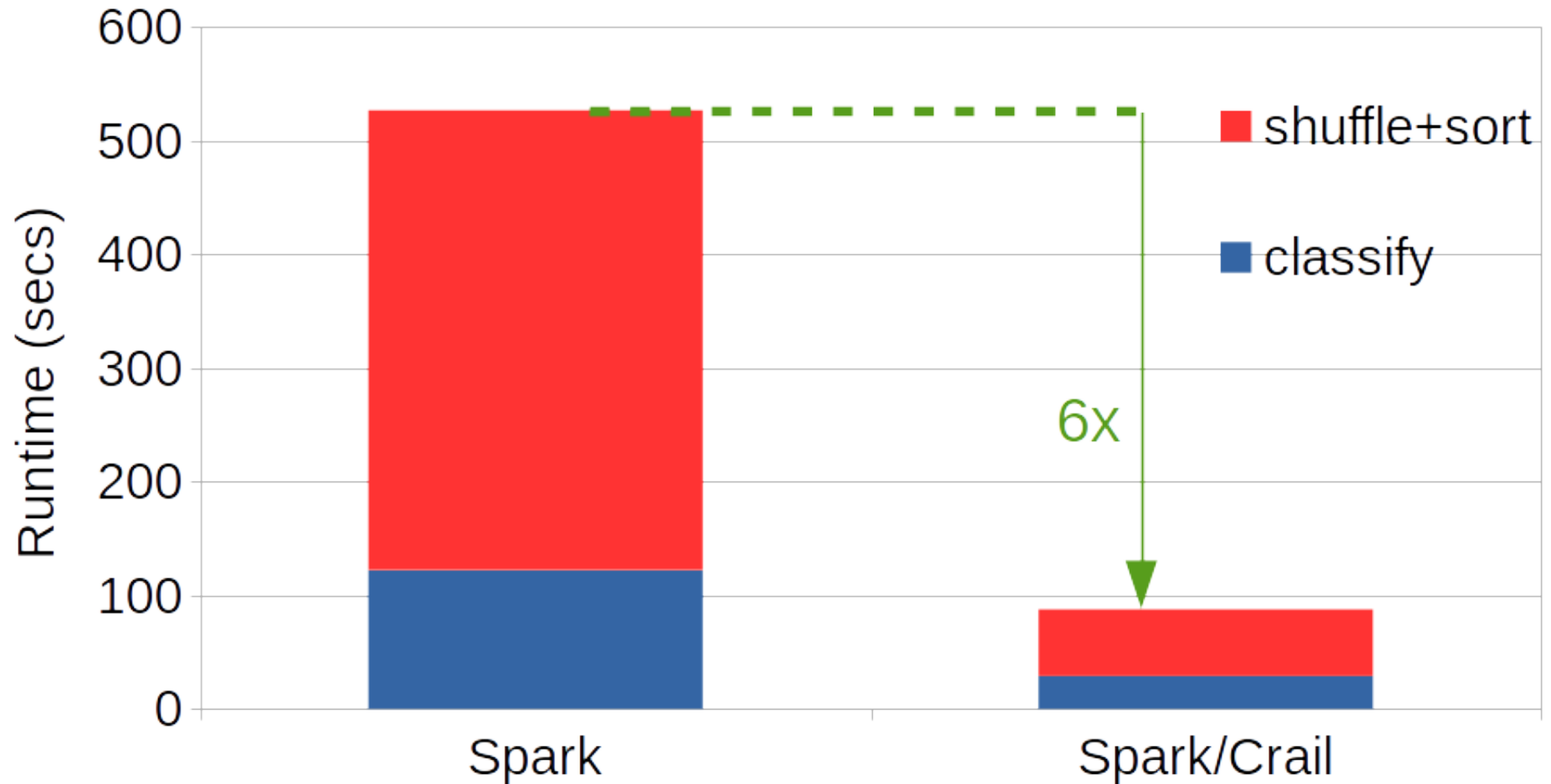
Crail offers 10-100x performance gains for Spark Broadcasts

Crail Workload - Spark - GroupBy



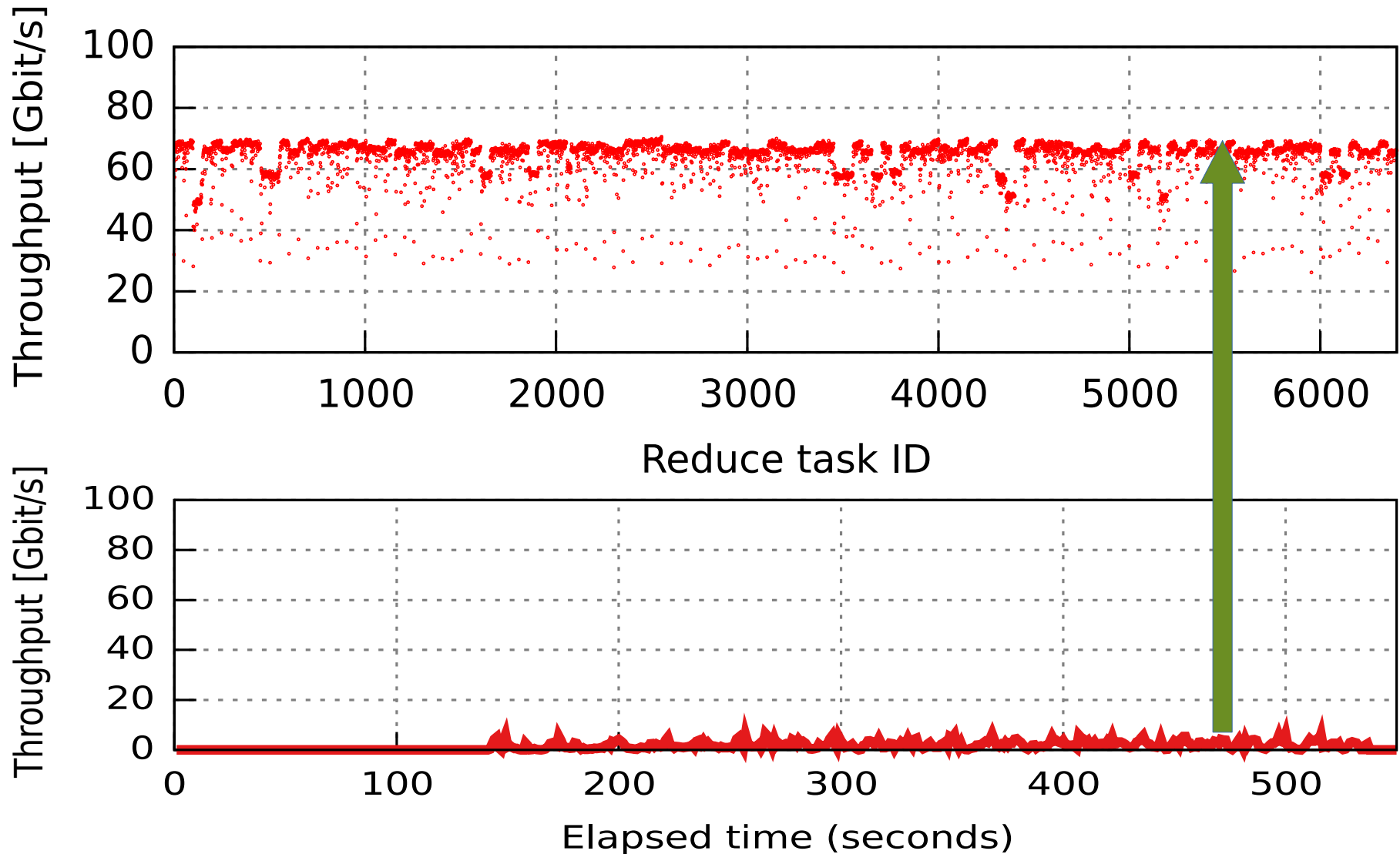
Crail offers 2-5x performance gains for Spark Shuffle

Crail Workload - Spark - TeraSort

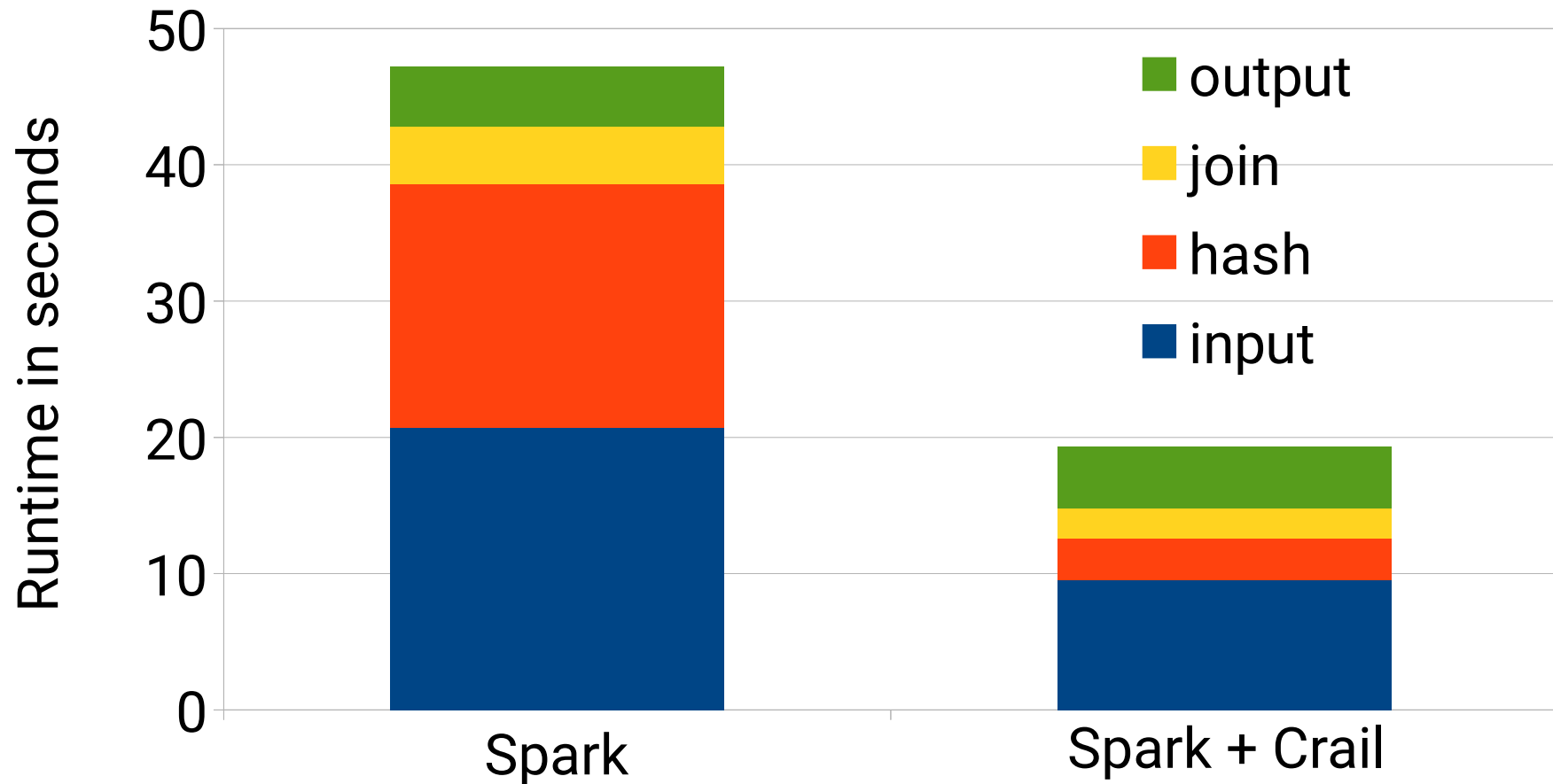


Crail offers 6x performance gains for Spark TeraSort [7]

Crail Workload - Spark - TeraSort

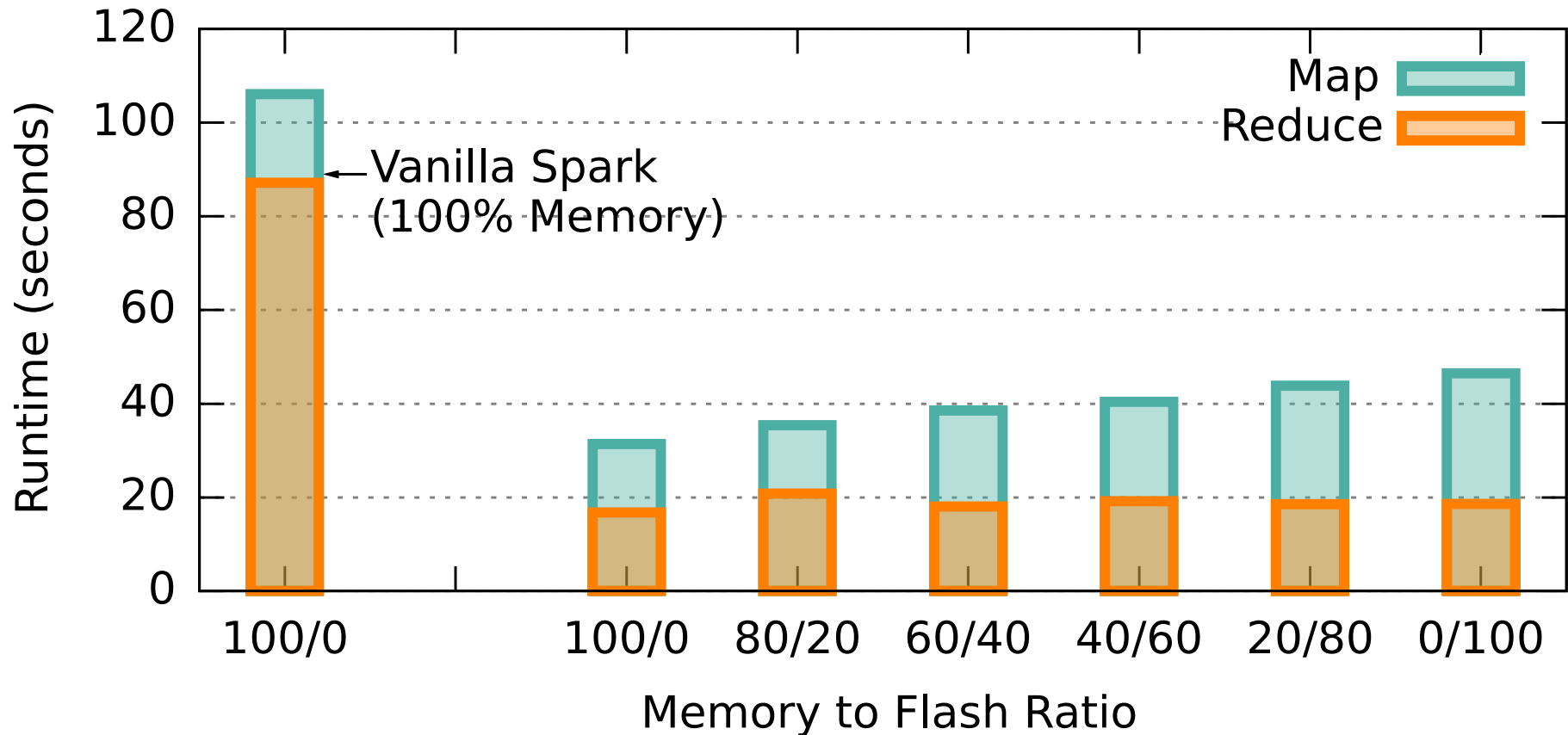


Crail Workload - Spark - EquiJoin



Crail offers 2x performance gains for SQL JOIN

Crail Store – Flash Disaggregation



Crail offers the possibility for storage dis-aggregation

Current Status and Future Plans

- Apache Incubator project since November, 2017
 - First source release (tag: 0db64fd), a few weeks ago
- Plenty of opportunities
 - Multiple languages support (C++ (WiP), Python, Rust, `_your_fav_language_`)
 - Multiple frameworks integration (Flink, Hadoop, λ compute, TensorFlow, SnapML[5], `_your_fav_framework_here_`)
 - Multiple datanode (network and storage) integration
 - Automated testing and packaging framework
 - Deploying in the cloud/containerization
 - JVM Optimizations
 - And all the fun stuff that you know and love about Apache projects !

Thanks to

Patrick Stuedi, Jonas Pfefferle, Michael Kaufmann,
Adrian Schuepbach, Bernard Metzler

Thank you !



See you all at the mailing list :)

- Website: <http://crail.incubator.apache.org/>
- Blog: <http://crail.incubator.apache.org/blog/>
- Mailings list: dev@crail.incubator.apache.org
- JIRA: <https://issues.apache.org/jira/browse/CRAIL>
- Slack: <https://the-asf.slack.com/messages/C8VDLDWMV>

References

- [1] DiSNI: Direct Storage and Networking Interface, <https://github.com/zrlio/disni>
- [2] jNVMf: A NVMe over Fabrics library for Java, <https://github.com/zrlio/jNVMf>
- [3] “Crail: A High-Performance I/O Architecture for Distributed Data Processing”, in the IEEE Bulletin of the Technical Committee on Data Engineering, Special Issue on Distributed Data Management with RDMA, Volume 40, pages 40-52, March, 2017.
<http://sites.computer.org/debull/A17mar/p38.pdf>
- [4] Crail I/O acceleration plugins for Apache Spark, <https://github.com/zrlio/crail-spark-io>
- [5] Snap machine learning (SnapML), <https://www.zurich.ibm.com/snapml/>
- [6] Apache Crail (Incubating) performance blogs
 - Part I: DRAM, <http://crail.incubator.apache.org/blog/2017/08/crail-memory.html>
 - Part II: NVMeF, <http://crail.incubator.apache.org/blog/2017/08/crail-nvme-fabrics-v1.html>
 - Part III: Metadata, <http://crail.incubator.apache.org/blog/2017/11/crail-metadata.html>
- [7] Sorting on a 100Gbit/s Cluster using Spark/Crail, <http://crail.incubator.apache.org/blog/2017/01/sorting.html>

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